

1 **Journal Title**

2 **A Practical Approach to Protect IoT Devices against Attacks and**  
3 **Compile Security Incident Datasets**

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15 **Abstract**

16 The Internet of Things (IoT) introduced the opportunity of remotely manipulating home appliances  
17 (such as heating systems, ovens, blinds, etc.) using computers and mobile devices. This idea  
18 fascinated people and originated a boom of IoT devices together with an increasing demand that  
19 was difficult to support. Many manufacturers quickly created hundreds of devices implementing  
20 functionalities, but neglected some critical issues pertaining to device security. This oversight gave  
21 rise to the current situation where thousands of devices remain unpatched having many security  
22 issues that manufacturers cannot address after the devices have been produced and deployed. This  
23 article presents our novel research protecting IOT devices using Berkeley Packet Filters (BPFs)  
24 and evaluates our findings with the aid of our Filter.tlk tool, which is able to facilitate the  
25 development of BPF expressions that can be executed by GNU/Linux systems with a low impact  
26 on network packet throughput.

27 **1. Introduction and motivation**

28 The evolution of Internet and communication networks from their emergence in the sixties to today  
29 has enabled a revolution in the way people and businesses interact. People today communicate  
30 worldwide using mobile devices, which have a reliable broadband (4G) Internet connection.  
31 Despite these great advances, Aceto *et al.* [1] note that network outages are still a challenge to  
32 solve because they are frequent, hard to fix, expensive and, in particular, poorly understood by

33 users. Whilst there exists a variety of problems surrounding network availability (Aceto *et al.* [1]),  
34 this study presents a proposal to avoid or at least minimize the effects of problems caused by  
35 software attacks through networks, including worms [2], [3] and remote attacks to exploit server  
36 vulnerabilities [4].

37  
38 Patching avoids the compromise of target systems through e.g. malware and vulnerability exploits.  
39 However, the growth of Internet of Things (IoT) applications running on devices, that frequently  
40 do not support patching, using Internet and TCP/IP networks for communication purposes limits  
41 this possibility. Moreover, software upgrades are not always immediately available when the  
42 vulnerability is discovered, as patch development and distribution depend on developer  
43 circumstances. The existence of proactive defense mechanisms [5] capable of mitigating risks  
44 associated with unpatched components within an otherwise trusted TCP/IP network would be very  
45 valuable. In this study, we take advantage of the firewall support of operating systems to develop  
46 a highly efficient mechanism to detect and bring together information about malicious network  
47 traffic.

48 Firewalling support has become an essential feature of modern operating systems. The use of  
49 firewalls is one of the easiest mechanisms to manage network defense. However, its effectiveness  
50 is clearly limited to protect IoT devices against malware and vulnerability exploiting. [6]. In the  
51 well-known Linux operating system, firewall capabilities have been provided primarily through  
52 packet filtering technology, and have evolved from a Netfilter ipfw system port (included in Linux  
53 kernel 1.1) to netfilter/iptables (included in Linux 2.4 kernel series). This evolution entailed the  
54 introduction of significant innovations such as the tracking of TCP connections or the possibility  
55 of altering packets in transit (mangle table). Despite the popularity of these filters, netfilter/iptables  
56 firewalling subsystem will be replaced in order to speed up filtering process and increase the  
57 information achieved for each packet to filter (such as payload information).

58 Wireshark [7] capture filters are defined by using libpcap filter language. Filter examples that are  
59 designed to detect some worms and exploits are available in Wireshark Wiki [8] showing the  
60 power of this filter syntax. The syntax of capture filters is commonly known as Berkeley Packet  
61 Filter (BPF) and is supported in the kernel of most UNIX-like operating systems. This syntax is  
62 also implemented by libpcap/Winpcap to be used at the user level in tools such as Wireshark. BPF  
63 [9] was first introduced in 1990 as a tool for capturing and filtering network packets that matched  
64 specific rules. BPF support was included in Linux Kernel by implementing a small virtual machine  
65 that runs compiled BPF programs injected from user-space [10]. Later, a BPF Just-In-Time (JIT)  
66 compiler was added to speed up the performance of the execution of bytecodes. Currently, BPF  
67 can be loaded for its execution into kernel with different tools to execute different tasks, such as  
68 system monitoring (through using *perf* tool), network traffic control and quality of service (through  
69 *tc* tool), and packet filtering (through *ip link* tool included in *iproute2* suite or *iptables*).

70 Due to its flexibility, BPF has been used by important technology companies such as Google,  
71 Facebook, Cloudflare and Netflix to address network security issues, load-balancing, traffic-  
72 filtering and monitoring [11]–[14]. A comparison of the filtering performance achieved by a BPF-  
73 based filter (BPFilter), iptables and nftables has also been provided in other studies [11], [15]  
74 showing that BPFilter runs up to 5 times faster than iptables. This scenario led to the consideration  
75 of BPF as a reliable candidate to replace iptables (and nftables) as the kernel firewall subsystem  
76 for Linux [11]. However, despite the fact that BPF syntax is more powerful than that offered by  
77 current Linux firewalls, BPFilter only takes advantage of the BPF virtual machine to speed up  
78 rules created by older tools. Majkowski [16], [17] demonstrated how to take advantage of BPF in  
79 conjunction with iptables to filter packages and define new chains. These works allowed system  
80 administrators to take advantage of the rich syntax and efficient execution of BPF expressions to  
81 filter packages in real environments and protect IoT devices against malware and vulnerability  
82 exploiting.

83 We developed Filter.tk to work in conjunction with these tools. Filter.tk is a framework to complete  
84 the full lifecycle (creation, debugging and testing) of BPF iptables-compliant pattern design for  
85 mitigating both worm and exploit attacks. The development of patterns will be useful for the future  
86 creation of a BPF rules database usable in the form of well-known community collaboration  
87 products such as Ansible Galaxy [18], [19] or DockerHub [20], where users can share BPF to  
88 protect IoT devices, computers and software against worm and exploit attacks. Additionally, the  
89 information about harmful network packets can be uploaded to centralized repository for research  
90 purposes. Particularly, this data, if compiled worldwide, could allow the identification of security  
91 threats and help in the identification of new offensive packet patterns.

92 The remainder of this paper is structured as follows: Section 2 introduces the state of the art in  
93 well-known worms, security vulnerabilities and IoT security. Section 3 introduces our proposal to  
94 address both the protection of devices against security vulnerabilities (and hence, worm attacks  
95 exploiting those vulnerabilities) and the compilation of security incident datasets in the context of  
96 IoT while Section 4 is centered on discovering the utility of the toolset through case studies.  
97 Finally, Section 5 presents the main conclusions of the work and outlines the directions for future  
98 research.

## 99 **2. State of art**

100 During the early 2000s, Internet Worms became very popular due to the effects of well-known  
101 worms such as Code Red (versions 1 and 2), Nimda, SQL Slammer or Blaster. Some of these  
102 worms are compiled in the work of Qing and Wen [2]. However, due to the heightened awareness  
103 of users and developers about the importance of security, this kind of malicious software is solely  
104 spread in P2P networks and operates in a passive form [21]. Instead of performing an active search  
105 to infect computers, passive worms require human intervention, i.e. by downloading an infected  
106 file from a P2P network, to replicate themselves. Despite the propagation of passive worms in P2P

107 networks mainly connected with the illegal downloading of software and multimedia materials,  
108 the dissemination of these Internet worms and their mitigation has been fairly well discussed in  
109 previous literature [21]–[29]. The detection of new vulnerabilities allowing remote exploitation is  
110 a very active area as evidenced by the latest exploits published in exploit-db [4]. Although the  
111 existence of vulnerabilities allowing the execution of remote commands could provide a  
112 mechanism for the dissemination of worms, the quick response of software development teams to  
113 provide security patches discourage malware developers from designing new worms. With this in  
114 mind, the goal of this work is to mitigate attacks exploiting software vulnerabilities, with a special  
115 interest in those targeting an IoT device.

116 Many IoT applications and devices have become available for smart home automation. Querying  
117 ‘remote’ ‘hardware’ exploits in exploit-db and other similar databases resulted in a number of  
118 exploitable vulnerabilities in well-known products (such as intelligent TVs, cameras, etc.). This  
119 shows that IoT developers have been prioritizing the development and creation of functionalities  
120 for most demanding users while frequently neglecting security considerations.

121 A few works have addressed issues in IoT security, such as the use of block-chain communications  
122 [30]–[33]. These usually refer to security issues pertaining to confidentiality, integrity, and  
123 availability in the communications between IoT devices and IoT. The work of Ammar *et al.* [34]  
124 provides a critical review of eight well known IoT frameworks with special emphasis on security  
125 issues (analyzing models and approaches provided for ensuring security and privacy, pros and cons  
126 of each framework in terms of fulfilling the security requirements and meeting the standard  
127 guidelines and identifying design flaws). Wood and Stankovic [35], [36] provide studies about  
128 network issues. Particularly, the former work is centered in security-related issues about IoT  
129 communication protocols whilst the later analyzes denial of service threats in IoT environments.

130 Wack *et al.* [37] review the risk of platform Software/Firmware vulnerabilities that enable the  
131 reception of malicious attacks. To the best of our knowledge, there is no research work focused on  
132 the prevention, management and response to vulnerability exploiting and worm attacks in IoT.  
133 Closing this gap, we studied how to take advantage of firewalling schemes to implement these  
134 protections.

## 135 **2.1. OS firewalling support**

136 Common OS firewalls, such as those that can be implemented through GNU/Linux kernel  
137 firewalling subsystem, are usually implemented as packet filters [37], [38], which consist of a  
138 default policy for packets and a sequence of rules that define the actions performed on packets  
139 when they satisfy certain conditions. Specifically, each firewalling rule contains a triggering  
140 condition, usually a simple condition or the logical AND of simple conditions, together with an  
141 action to execute when the rule is triggered. Triggering conditions are defined over the second,  
142 third and fourth TCP/IP layers. The support for stateful inspection of connections is available for  
143 kernel versions 2.4 and above [39].

144 The first GNU/Linux firewall generation was included on 1.1 kernel through an implementation  
145 of *ipfw* functionalities contributed by Alan Cox [40]. The *ipfwadm* user-space tool was used to  
146 configure the *ipfw* services offered by the kernel [41]. These kernels allowed defining three  
147 different firewall filters to handle (i) input packets (-I ipfwadm argument), (ii) output packets (-O  
148 and (iii) forwarded packets (-F, used in conjunction with *ip\_forwarding* feature). *Accept*, *deny*  
149 (discard the packet) and *reject* actions were used either for rules (-a command parameter) or as  
150 default policy (-p). In order to create and design the trigger condition of each rule, the system  
151 administrator can test for the protocol (TCP, UDP, ICMP or IP), the port (for TCP and UDP) or  
152 the ICMP type. Logging is supported through the -o modifier.

153 The support for ipfw was replaced by *ipchains* in the 2.2 version of the kernel [42]. One of the  
154 most important changes in the *ipchains* scheme was the introduction of *chains* to help reduce the  
155 computational cost and facilitate its design [43]. As opposed to the others, *ipchains* firewalls  
156 include INPUT, OUTPUT and FORWARD chains, which bring together filtering rules applied to  
157 packets where the current computer is the destination, the origin or a router for the packet  
158 respectively. Each firewall *chain* is composed of a ruleset and a default policy. The default policy  
159 is applied to packets that do not match any rule. The existence of default policies allows defining  
160 firewalls using two different schemes: (i) accept all except those packets explicitly denied or (ii)  
161 deny all except those packets explicitly accepted. Of these, the latter is advisable for security  
162 reasons.

163 New functionalities offered by ipchains with regard to ipfw were quite limited so it was quickly  
164 replaced by iptables (in Linux 2.4 series) [44]. Iptables/netfilter included the table concept to bring  
165 together chains with similarities. Iptables included the firewall tables *filter*, *nat*, and *mangle*. The  
166 first one included three chains to support the filtering of input, output and forwarded connections  
167 (INPUT, OUTPUT and FORWARD respectively). The NAT table is composed of PREROUTING  
168 and POSTROUTING *chains* to add rules to change destination or source addresses respectively.  
169 To this end, rules included in these chains could only use DNAT, SNAT, REDIRECT or  
170 MASQUERADE actions. Finally, MANGLE table allows marking packages for further processing  
171 and modifying some parameters of packets including TOS or TTL. These actions could be  
172 executed by using MARK, TOS and TTL actions. Finally, iptables brought the stateful packet  
173 inspection to Linux firewalls making it possible to determine whether a packet belongs to an  
174 established TCP connection (-m state --state=ESTABLISHED) or, conversely, is connected with  
175 other previous packets (-m state --state=RELATED).

176 Iptables have been widely used to implement packet filters on Linux for many years [45].  
177 However, some limitations of iptables, such as the existence of a unique action for a rule or the  
178 complexity of the syntax, led to the creation of other filtering frameworks. Hence, nftables  
179 emerged as an iptables replacement on kernel version 3.13 (2013) [44]. Nftables included a  
180 completely new and fresh syntax that avoided the need to use hyphens and the uppercase/lowercase

181 flags. The use of nftables allows tables and chains to be created with specific names and associated  
182 with hooks, thus avoiding the strict tables/chains structure defined by iptables.

183 Despite the new functionalities of nftables, most Linux users continue using the old iptables  
184 framework, in part due to the numerous changes in the syntax which hindered adoption. Some  
185 translation utilities were introduced to aid in the migration from iptables to nftables [46]. In  
186 addition, nftables work as a sequential filter whereby every packet is matched one by one against  
187 a list of rules. The speed of checking the rules is quite limited (up to three times slower than using  
188 BPF) [11], which led to the emergence of bpfILTER as a new Linux firewalling subsystem [47] able  
189 to outperform the speed of previous filtering alternatives [11]. BpfILTER has been added  
190 experimentally to Linux Kernel 3.18, now allowing nftables and iptables rules to be executed by  
191 Linux kernel as BPF.

192 Standard definitions within BPF only allow current packet filtering firewall [48] schemes to  
193 analyze some information from packet headers (such as IP and MAC addresses, ports, TCP flags,  
194 ICMP types, etc.) and packet state. Additional new features would improve the firewalling  
195 performance, such as analyzing the payload of the packet or the information about application  
196 layer protocols. These features are frequently included in deep packet inspection techniques [49],  
197 [50] but are often too slow to be included in standard firewalls. In order to provide a deep  
198 description of packets on the firewall layer and quickly evaluate them, the use of BPF language  
199 together with the BPF virtual machine subsystem included in the current versions of Linux kernel  
200 seems to be an elegant solution, especially as BPF had been used before to accomplish similar  
201 difficult network tasks with low computational effort [11].

202 Despite its performance and low computational costs, developing a BPF-based firewall able to  
203 exploit full packet data (headers and payload) is a hard task that would require both the existence  
204 of tools to aid in the development of conditions, and packet datasets. CapLoader [51] and  
205 Wireshark/tcpdump [7], which can also be integrated with NDPI [52], are capable of loading and  
206 analyzing packets included in PCAP files and check whether a BPF expression match them. These  
207 tools can be successfully executed with large collections of packets, such as the shared by Netresec  
208 [53]. However, the design of BPF filters is not easy and should be simplified to impact on real-  
209 world firewall applications. Similarly, the evaluation of BPF filters should be automated to  
210 improve performance. Both the simplification and automation have been addressed by our  
211 Filter.tlk toolset and are the main contribution of this work. Both Filter.lk functionality and its  
212 practical use is described in the next section.

### 213 **3. Filter.tlk**

214 This section provides a comprehensive description of the design architecture of Filter.tlk [54] tool  
215 and documents the process of creating customized filters to classify network traffic according to  
216 the content of the packets.

217 Filter.tlk comprises three different utilities to aid in the creation of BPF filters: (i) an interface to  
 218 design BPF filtering conditions, (ii) a Wireshark LUA plugin to automate the testing of BPF filters  
 219 with PCAP packet datasets and (iii) a script to easily compile BPF filters and create iptables rules.  
 220 Figure 1 shows the different components included in Filter.tlk and their use in a real environment.

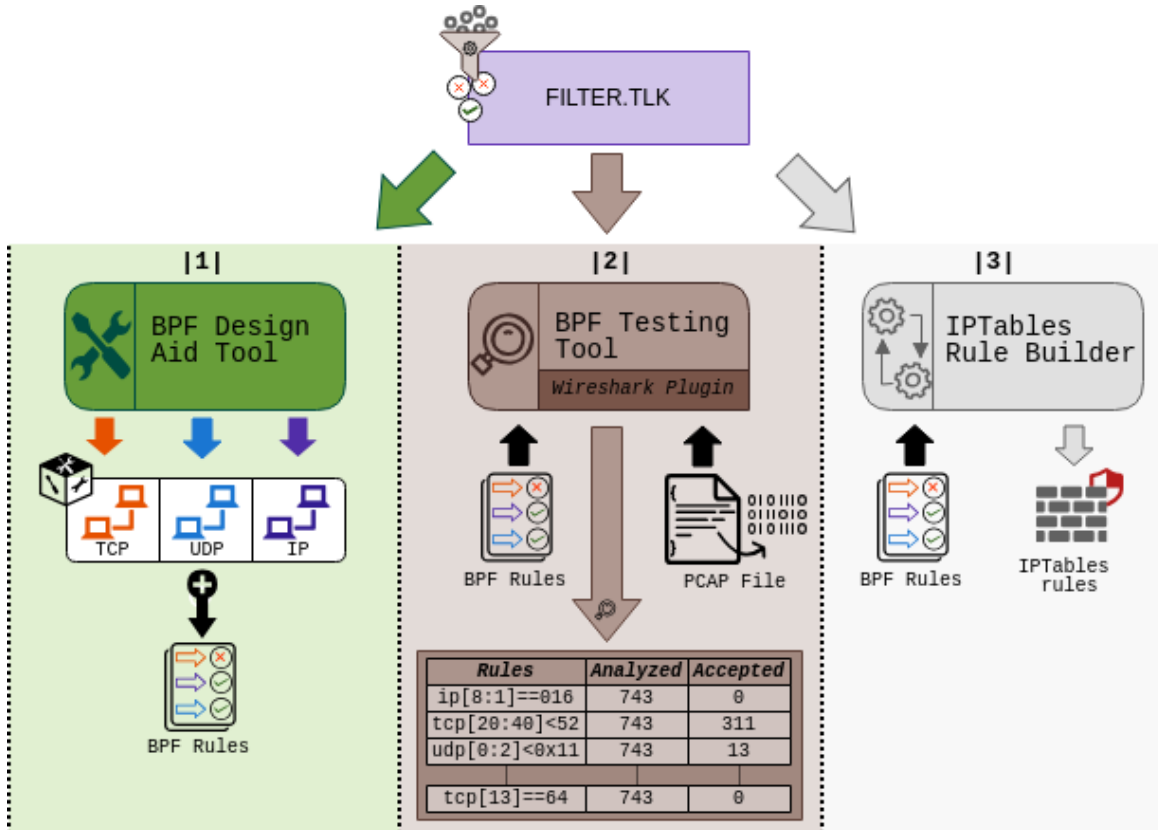


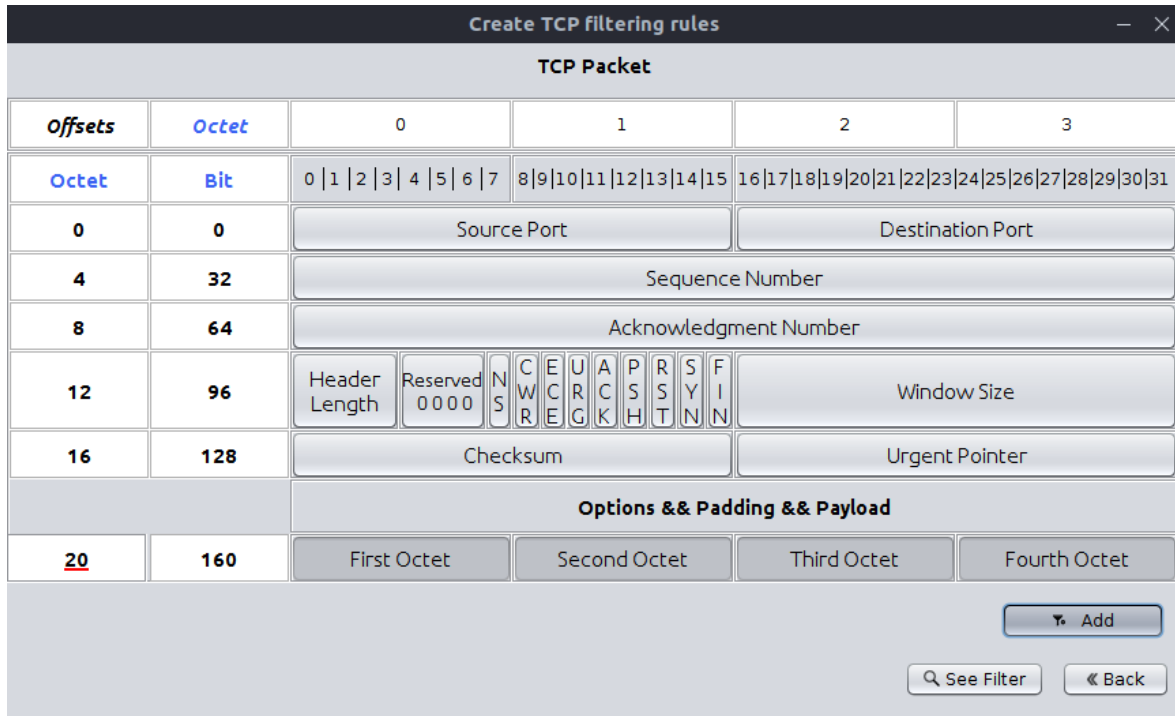
Figure 1. Filter.tlk architecture

221  
 222

223 As we can deduct from Figure 1, the design of a firewall rule with Filter.tlk comprises three stages  
 224 that are made with different tools included in the package. We begin by taking advantage of the  
 225 BDAT (*BPF Design Aid Tool*) to design a filtering condition to detect a certain kind of packet.  
 226 The designed BPF filters can then be tested with different packet sets (a set of packets matching  
 227 the pcap filter and others mismatching the pcap filter) using BPF Testing tool (BTT). BTT is able  
 228 to easily assess the quality of an input filter by using different datasets. Once BPF rules have been  
 229 tested, they can be easily transformed into iptables rules using IPTables Rule Builder (IPTRB)  
 230 script.

231 BDAT is responsible for creating BPF filters as conditions defined from transport and network  
 232 layers (see Figure 1). By using BDAT through a simple graphical interface, we can create a  
 233 Boolean BPF filter evaluating expressions related to network or transport headers and payloads  
 234 (UDP, TCP, IP, ICMP). In order to create header conditions, users must select the field of the  
 235 header on which they want to establish the condition that the filter must fulfill. Once the condition  
 236 has been defined, the user can continue adding new conditions for the same filter or create a new

237 one. Once all filters are defined, they can be exported to a file for testing in BTT (*BPF Testing*  
 238 *Tool*). As an example, Figure 2 shows how administrators can easily incorporate a condition about  
 239 a HTTP POST request by specifying conditions about TCP payload.



a) Select TCP payload octets



b) Establish values for octets



240

Figure 2. BPF rule definition process.

241 As depicted in Figure 2a, we selected the first four octets from TCP payload for comparison  
242 purposes. BDAT allows selecting one, two or four octets from each word (32bits) to check the  
243 condition. The next step of the wizard (see Figure 2b) allows to easily define the value using  
244 Hexadecimal, ASCII or Decimal notations. In order to compare any octet from payload (and  
245 options/padding), the offset value (highlighted in red in Figure 2a) can be edited to the desired  
246 value. Please note that the designed condition is provided as example and should be complemented  
247 with a “header length” value of five to ensure the absence of options/padding field. In the next step  
248 of the wizard, the current BPF condition rule is added to the whole BPF filter, allowing the  
249 generation of filters comprising multiple tests.

250 BTT is a plugin for Wireshark that applies a filter or set of filters over a pcap file. As a result, we  
251 obtain information for each applied filter -about the number of packages analyzed, accepted and  
252 rejected. A set of pcaps with the packets accepted by the filter grouped by the destination IP is also  
253 provided for debugging purposes. In order to detect errors, each rule should be tested using a pcap  
254 database containing only the packets that should be captured (ensuring a result of 0 rejected is  
255 achieved) and another one containing normal network traffic (guaranteeing a result of 0 accepted  
256 is achieved).

257 Finally, IPTRB (*IPTables Rule Builder*) script can transform the BPF uncompiled filters into full  
258 featured IPTables rules. The process is guided by an intuitive libncurses-based graphical user  
259 interface that allows customizing the generated rule. The rule can be generated for filtering and/or  
260 harmful packet logging purposes. A scheduled task (i.e. crontab) could be periodically executed  
261 (for instance once a day) to upload the compiled information (logs) to a centralized repository for  
262 its further analysis.

### 263 **3.1. Filter.tlk implementation**

264 This subsection provides a brief description of the most relevant implementation issues for the  
265 development of each tool included in Filter.tlk.

266 BDAT was designed as a Java standalone application, which can easily be executed using any Java  
267 Virtual Machine implementation. The interface was designed using JFC/Swing library [55].

268 BTT is a Wireshark plugin that was implemented using the Lua programming language [56], which  
269 is supported by Wireshark for the development of new functionalities, such as the creation of  
270 dissectors or listeners [57]. The dissectors are intended to analyze part of the data of a packet,  
271 while the listeners are used to count the number of occurrences of an event; for example, the  
272 number of packets matching a filter. In this study, we used Lua language to implement a Wireshark  
273 listener to evaluate filters and count packets fitting the target BPF condition (accepted) or not  
274 (rejected).

275 IPTRB is a bash script that combines the use of dialog command [58] to provide an easy-to-use  
276 intuitive graphical user interface. Moreover, the compilation of BPF rules into bytecode is done  
277 by using tcpdump [59] functionalities.

278 Finally, we used Ansible [19] (a well-known IT Automation tool) to automate the installation of  
279 Filter.tlk in all supported platforms. Ansible is a popular IT automation tool whose main features  
280 are: (i) avoiding the need of scripts and/or custom code to deploy and update applications, and (ii)  
281 replacing agents on remote systems by standard SSH tools. The installation script was provided  
282 for Debian-based GNU Linux distributions.

### 283 **3.2. Deployment of generated filters**

284 To take advantage of expressions (iptables rules and BPF) generated using our BPF framework,  
285 we consider two different scenarios: (i) IoT devices using a GNU Linux-based software/firmware  
286 (ii) and other IoT devices with no BPF/iptables support. In the first scenario, iptables rules can be  
287 directly integrated into the firmware to protect them against malicious attacks. We are working on  
288 the development of a service to share BPFs together with a tool able to automatically download  
289 and upgrade BPFs for different IoT devices.

290 Although GNU/Linux is present in some devices, there are many appliances running other  
291 Operating Systems where the execution of BPF is not possible. Taking this into account, we are  
292 working on the design of a small bridge router (brouter) [60] device running GNU/Linux and  
293 ebttables. A brouter is a device that is able to transparently forward all traffic between two ethernet  
294 interfaces, and that allows the inclusion of filtering rules for network interfaces. This solution  
295 would be applicable for IoT devices connected to the network through an ethernet connection.

296 The main weakness of using BPF filters to protect devices against attacks is that we are unable to  
297 protect 802.11-based (WLAN, Wireless Local Area Network) IoT devices that do not run  
298 GNU/Linux.

299 Next section presents a comprehensive practical example describing the process of using the  
300 Filter.tlk tool to design a filter capable of detecting and filtering two important vulnerabilities  
301 recently discovered on well-known IoT devices.

## 302 **4. Experiments**

303 In this section, we test the filter designed to detect and filter attacks using two vulnerabilities in  
304 two well-known IoT devices that allow the remote execution of arbitrary commands: (i) LG  
305 Supersign TVs, and (ii) ASUS ADSL Router DSL-N12E\_C1. The next subsection shows the work  
306 environment prepared in order to generate high quality BPF patterns. Moreover, subsection 4.2  
307 presents the experimental protocol and results of our case studies while subsection 4.3 measures  
308 the impact of the use of these filters in the performance of IoT devices. Finally, subsection 4.4

309 shows how to compile and take advantage of the information gathered by IoT devices using  
310 Filter.tlk for scientific purposes.

#### 311 4.1. Configuring the working environment

312 In order to generate high quality BPF expressions that describe the pattern of a vulnerability  
313 exploitation, the use of a large packet database for testing purposes is advisable. Fortunately, many  
314 publicly available packet datasets can be freely downloaded from the Internet. Table 1 compiles a  
315 list of useful datasets.

316 Table 1: Publicly available datasets

Dataset	Number of packets	Number of files	Format	Short description
contagiodump[61]	988898	1154	pcap zipped	Collect malicious and exploit pcaps from various public resources (2013-2015)
Malware Traffic Analysis [62]	2445211	1291	pcap zipped	Malicious network traffic (2013-2018)
GTISK PANDA Malrec [63]	100201	373	pcap	Malware samples run in PANDA (2018)

317 From the datasets included in Table 1 and other sources, we built up a group of packets that would  
318 be used to ensure that inoffensive network requests are not captured by the designed BPF  
319 expression.

320 We decided to study and generate BPF filters for two vulnerabilities of well-known IoT devices.  
321 In order to determine the quality of the BPF expressions created using a BTT Lua script, we used  
322 as negative samples the conjunction of all packets from sources introduced in Table 1 and other  
323 legitimate packet sets compiled by us. In order to aggregate all negative samples in a single pcap  
324 file, we combined all sources using the mergecap [64] tool provided by Wireshark.

#### 325 4.2. Executing the experiment

326 Recently, a vulnerability allowing remote execution of arbitrary commands appeared on LG  
327 SuperSign TVs (CVE-2018-17173) [65], [66]. These smart TVs include a CMS running on the top  
328 of LG webOS 3.3 (a Linux-based OS). The discovered vulnerability allows remote code execution  
329 (by achieving a reverse shell connection) by taking advantage of the URL used to see thumbnails  
330 of the user images. We used the exploit versions to generate a pcap file capturing the attacks. The  
331 Filter.tk comprises a three stages operation. The first step designs the filtering condition to detect  
332 the packets, the second step tests the BPF filter with a set of packets captured in a pcap filter and  
333 with a set of normal packets. The third step converts the BPF into iptables rules. The generated  
334 BPF expression is shown in Table 2. These BPF expressions could be directly included in LG  
335 webOS to protect the TV.

336 Table 2. BPF expression to mitigate CVE-2018-17173 vulnerability

BPF expression	ip[2:2] > 0x008A and ip[9]==0x06 and tcp[2:2]==0x2378 and tcp[32]==0x47 and tcp[77:4]==0x3d253237 and tcp[81:4]==0x2532302d and tcp[85]==0x3b		
BPF assembler code	Bytecode		
(000) ldh [12]	22		
(001) jeq #0x800 jt 2 jf 21	40 0 0 12		
(002) ldh [16]	21 0 19 2048		
(003) jgt #0x8a jt 4 jf 21	40 0 0 16		
(004) ldb [23]	37 0 17 138		
(005) jeq #0x6 jt 6 jf 21	48 0 0 23		
(006) jeq #0x6 jt 7 jf 21	21 0 15 6		
(007) ldh [20]	21 0 14 6		
(008) jset #0x1fff jt 21 jf 9	40 0 0 20		
(009) ldxb 4*([14]&0xf)	69 12 0 8191		
(010) ldh [x + 16]	177 0 0 14		
(011) jeq #0x2378 jt 12 jf 21	72 0 0 16		
(012) ldb [x + 46]	21 0 9 9080		
(013) jeq #0x47 jt 14 jf 21	80 0 0 46		
(014) ld [x + 91]	21 0 7 71		
(015) jeq #0x3d253237 jt 16 jf 21	64 0 0 91		
(016) ld [x + 95]	21 0 5 1025847863		
(017) jeq #0x2532302d jt 18 jf 21	64 0 0 95		
(018) ldb [x + 99]	21 0 3 624046125		
(019) jeq #0x3b jt 20 jf 21	80 0 0 99		
(020) ret #262144	21 0 1 59		
(021) ret #0	6 0 0 262144		
	6 0 0 0		
Iptables commands	<pre>iptables -t filter -A INPUT -m bpf --bytecode "22,40 0 0 12,21 0 19 2048,40 0 0 16,37 0 17 138,48 0 0 23,21 0 15 6,21 0 14 6,40 0 0 20,69 12 0 8191,177 0 0 14,72 0 0 16,21 0 9 9080,80 0 0 46,21 0 7 71,64 0 0 91,21 0 5 1025847863,64 0 0 95,21 0 3 624046125,80 0 0 99,21 0 1 59,6 0 0 262144,6 0 0 0" -j DROP iptables -t filter -A INPUT -m bpf --bytecode "22,40 0 0 12,21 0 19 2048,40 0 0 16,37 0 17 138,48 0 0 23,21 0 15 6,21 0 14 6,40 0 0 20,69 12 0 8191,177 0 0 14,72 0 0 16,21 0 9 9080,80 0 0 46,21 0 7 71,64 0 0 91,21 0 5 1025847863,64 0 0 95,21 0 3 624046125,80 0 0 99,21 0 1 59,6 0 0 262144,6 0 0 0" -j LOG --log-prefix "Filter.tlk"</pre>		

337 The second analysis is about a remote code execution vulnerability in ASUS DSL-N12E\_C1  
338 router, specifically in firmware version 1.1.2.3\_345 (CVE-2018-15887) [67]. This vulnerability  
339 has been classified as critical because it allows the execution of arbitrary code using an unknown  
340 function of the file ‘Main\_Analysis\_Content.asp’. A remote attacker can then access the router as  
341 a privileged user via telnet application and run OS commands. Again, we take advantage of our  
342 framework to generate BPF expressions to filter this type of attack. The generated BPF is shown  
343 in Table 3.

344 Table 3. BPF expression to mitigate CVE-2018-15887 vulnerability

BPF expression	ip[2:2] > 0x0174 and ip[9]== 0x06 and tcp[2:2]== 0x0050 and tcp[32]==0x47 and tcp[326:4]==0x3d253630		
BPF assembler code	Bytecode		
(000) ldh [12]	18		
(001) jeq #0x800 jt 2 jf 17	40 0 0 12		
(002) ldh [16]	21 0 15 2048		
(003) jgt #0x174 jt 4 jf 17	40 0 0 16		
(004) ldb [23]	37 0 13 372		
(005) jeq #0x6 jt 6 jf 17	48 0 0 23		
(006) jeq #0x6 jt 7 jf 17	21 0 11 6		
(007) ldh [20]	21 0 10 6		
(008) jset #0x1fff jt 17 jf 9	40 0 0 20		
(009) ldxb 4*([14]&0xf)	69 8 0 8191		
(010) ldh [x + 16]	177 0 0 14		
(011) jeq #0x50 jt 12 jf 17	72 0 0 16		
(012) ldb [x + 46]	21 0 5 80		
(013) jeq #0x47 jt 14 jf 17	80 0 0 46		
(014) ld [x + 340]	21 0 3 71		
(015) jeq #0x3d253630 jt 16 jf 17	64 0 0 340		
(016) ret #262144	21 0 1 1025848880		
(017) ret #0	6 0 0 262144 6 0 0 0		
Iptables commands	<pre>iptables -t filter -A INPUT -m bpf --bytecode "18,40 0 0 12,21 0 15 2048,40 0 0 16,37 0 13 372,48 0 0 23,21 0 11 6,21 0 10 6,40 0 0 20,69 8 0 8191,177 0 0 14,72 0 0 16,21 0 5 80,80 0 0 46,21 0 3 71,64 0 0 340,21 0 1 1025848880,6 0 0 262144,6 0 0 0" -j DROP iptables -t filter -A INPUT -m bpf --bytecode "18,40 0 0 12,21 0 15 2048,40 0 0 16,37 0 13 372,48 0 0 23,21 0 11 6,21 0 10 6,40 0 0 20,69 8 0 8191,177 0 0 14,72 0 0 16,21 0 5 80,80 0 0 46,21 0 3 71,64 0 0 340,21 0 1 1025848880,6 0 0 262144,6 0 0 0" -j LOG --log-prefix "Filter.tlk"</pre>		

345 As shown in Tables 2 and 3, an iptables command can be easily generated from a BPF expression  
346 to drop and log packets that match it. Although iptables can only check expressions in the header  
347 of network packets, BPF expressions make it possible to examine information included in both  
348 packet headers and payload in order to find any potential exploitation of vulnerabilities. The  
349 second generated iptables rule allows for storing security information that can be uploaded to a  
350 centralized repository for its further analysis. In next section, we evaluate the performance impact  
351 on the IoT devices when using these filters.

### 352 4.3. The impact on filter throughput

353 We assessed the impact of using BPF filters on IoT devices in order to determine if they could be  
354 successfully used to protect IoT devices against network vulnerability exploitation. To perform  
355 this analysis we used an Apache web server installed on a Raspberry Pi 2 Model B [68].

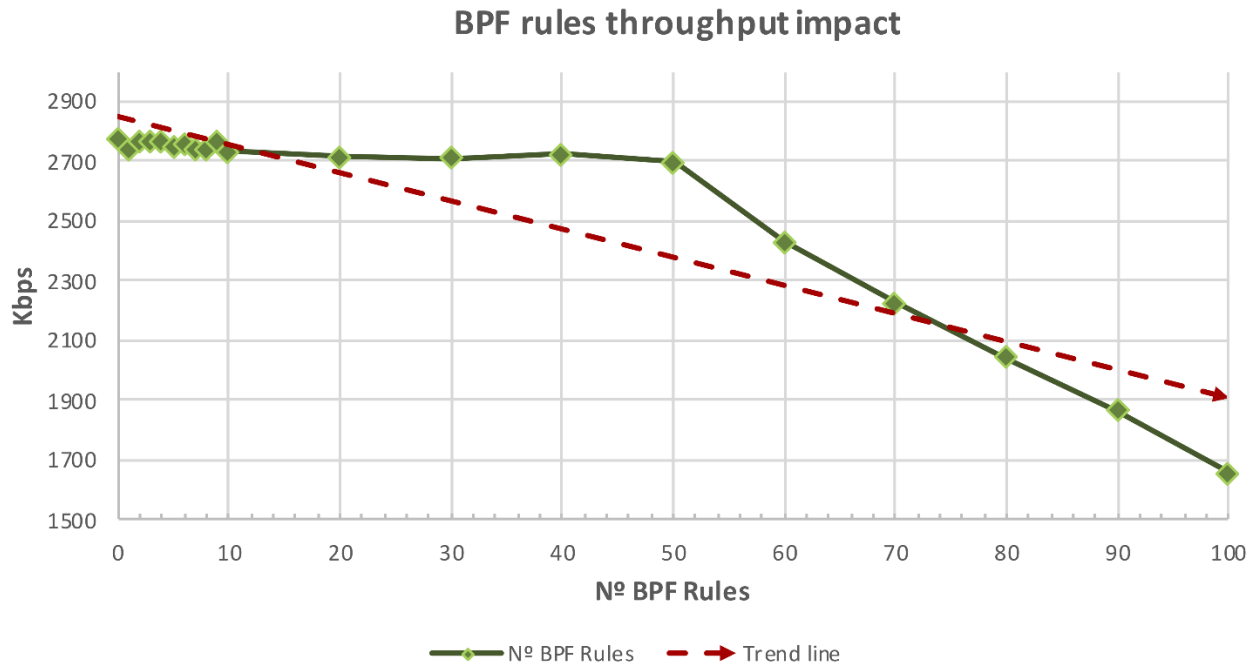
356 We leveraged the functionalities of Apache HTTP server benchmarking [69] and GNU parallel  
357 [70] tools to evaluate the impact of using BPF firewalls in IoT hardware. Using these tools, we  
358 benchmarked the execution of two parallel tests making 10000 HTTP requests distributed in 10  
359 threads, with 1000 requests per thread. The average of measurements made for parallelized tests  
360 is provided as result. For comparison purposes, we used the generated BPF expressions for the two  
361 case studies shown before. Table 4 compares the performance between the absence of attack  
362 protections and the usage of two BPF filtering rules.

363

Table 4. Performance impact when using BPF filters

	No protection	With BPF (2 rules)
HTML transferred (bytes)	107010000	107010000
Concurrent Time per request (ms)	1.9345	1.937
Time per request (ms)	19.3465	19.367
Time taken for tests (seconds)	19.3465	19.367
Total transferred (bytes)	109750000	109750000
Transfer rate (Kbytes/sec)	5539.905	5534.1

365 The results compiled in Table 4 show that the performance impact when using BPF filters is quite  
 366 limited and will not severely affect the overall operation of IoT devices. We analyzed the impact  
 367 of progressively adding BPF rules to the filter by adding up to 100 new rules and measured the  
 368 transfer rate after each BPF expression was added (see Figure 3). As long as the performance is  
 369 highly influenced by the presence of additional traffic in network and other processes consuming  
 370 CPU, we plotted a trend line to observe the degradation.



371

372

Figure 3. Analysis of the impact of BPF rules in throughput

373 As can be seen from Figure 3, the throughput degradation is close to zero when using up to 50  
 374 (non-fitting) BPF rules. However, the inclusion of more than 50 rules clearly damages the  
 375 performance of GNU/Linux firewalling system and would require the usage of additional iptables  
 376 speedup strategies, such as the creation of additional chains [71] and counters-based optimizations  
 377 [72].

378

379

One of the most interesting features of Filter.tlk framework is the compilation of information about worldwide IoT security incidents. The information gathered could be successfully analyzed using

380 Machine Learning techniques to provide worthwhile knowledge about: (i) the origin of the threat,  
381 (ii) better patterns for traffic filtering or (iii) the threat scale. Studying information about systems  
382 from which the attack is performed, we can successfully identify worms exploiting a certain  
383 vulnerability, the presence of an individual hacker, the execution of Distributed Denial of Service  
384 attacks or botnets.

385 Offering IoT users a product to protect their devices against attacks, whilst at the same time  
386 achieving information about dangerous offensive network packets targeting IoT products, will  
387 replicate a threat response model undertaken by traditional antivirus products. This knowledge  
388 allows the identification of better and perhaps simpler BPF patterns that can be used for network  
389 intrusion detection.

## 390 **5. Conclusions and future work**

391 In this paper, we have introduced an easy-to-use framework designed to aid in the development of  
392 fast firewalls based on using BPF, which can be executed by using standard firewall capabilities  
393 included in the Linux kernel (IPTables/Netfilter). These firewalls have been specifically conceived  
394 to protect IoT devices against the exploitation of remote vulnerabilities. Since the use of BPF  
395 bytecode can drastically speed up the execution of firewalls, we designed a collection of tools to  
396 facilitate the inclusion of BPF into firewalling rules. An experiment was carried out for the  
397 application of the introduced toolset.

398 Since BPF is one of the most efficient forms of filtering traffic, it provides a reliable solution for  
399 filtering in the context of IoT. Although, the use of specific BPF filters allows using payload  
400 information included in packets, it can only be directly implemented in devices using a GNU/Linux  
401 based firmware. We are currently working on the design of specific hardware to overcome the  
402 limitations of non-GNU/Linux Ethernet IoT devices and on the development of a package manager  
403 to automatically download BPF filtering strings and configure the firewall.

404 One of the most relevant functionalities of this scheme is the ability to easily build a dataset with  
405 the security incidents occurred in worldwide IoT devices, such as VizSec [73]. Future work will  
406 include mechanisms for analyzing them to achieve valuable security knowledge. We consider  
407 evolutionary computation as a candidate method for automatic filter generation through packet  
408 captures and consider DPI (Deep Packet Inspection [52]) to be a reliable way of simplifying  
409 filtering conditions, since it allows access to Application Layer information to define matching  
410 expressions. While DPI expressions cannot be directly included in BPF filters, we believe that they  
411 could be automatically transformed into simple BPF expressions to simplify the generation of BPF  
412 filters.

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